

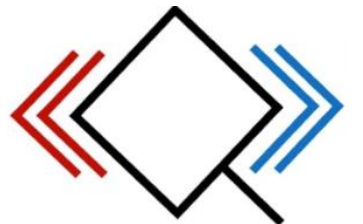
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Rhea Alexander

*npj Quantum Inf* **12**, 15 (2026)



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*Part 1: What and why?*

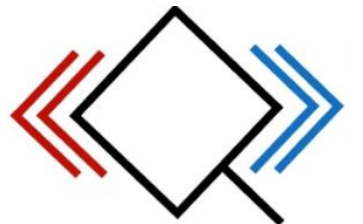
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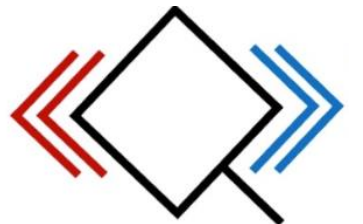
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*Part 2: Main results*

*Part 3: Methods: QEC + transversal gates = asymmetry distillation*

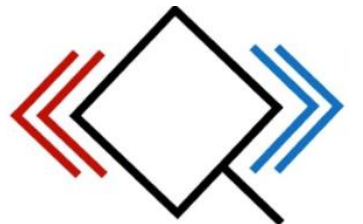
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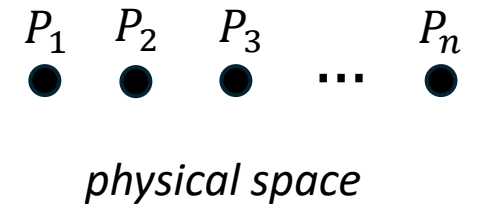
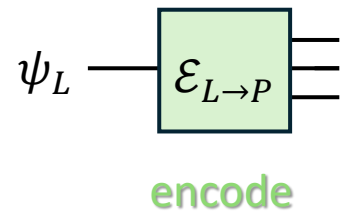
# The goal...



*\*fault-tolerant quantum computer\**

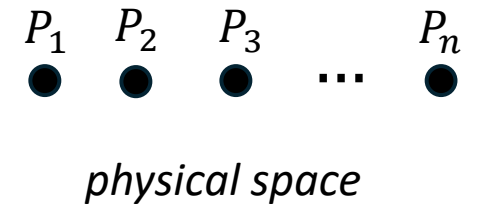
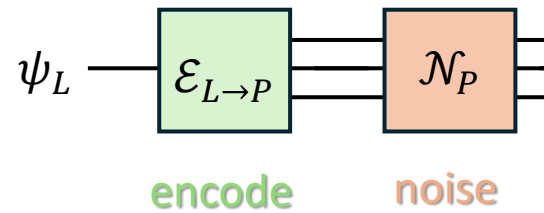
# Ingredients for fault-tolerance

## 1. “Error resistant” quantum memory: QEC code



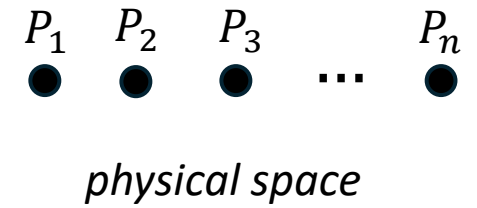
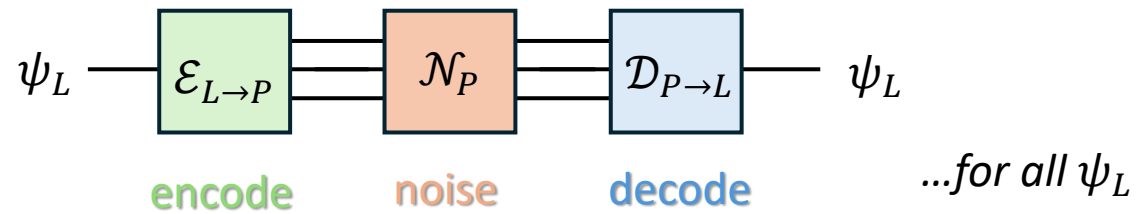
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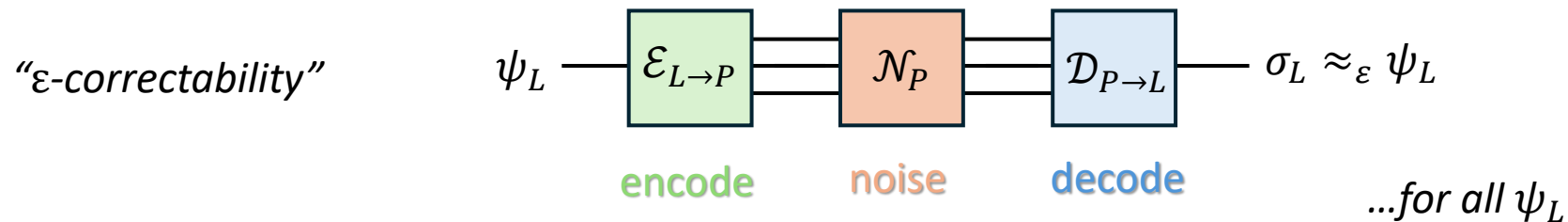
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- Any fault-tolerant architecture must circumvent!
- **Approximate QEC:** relax assumption of *exact* error correction<sup>2</sup>



# Approximate Eastin-Knill

- “Robust” variants of Eastin-Knill provide lower bounds on asymptotic QEC inaccuracy of codes with UTG (necessary conditions)[1]

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**Q:** *necessary and sufficient conditions for single-shot regime?*

## Part 2: Main results

# Main result

**Theorem.** Any encoder  $\mathcal{E}_{L \rightarrow P}$  with a universal set of transversal gates is  $\varepsilon$ -correctable w.r.t erasure of  $m$  subsystems  $\mathcal{N}_{P \rightarrow P'} := \text{tr}_{P_1 \dots P_m}$  iff

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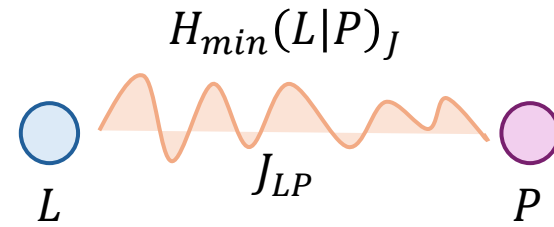
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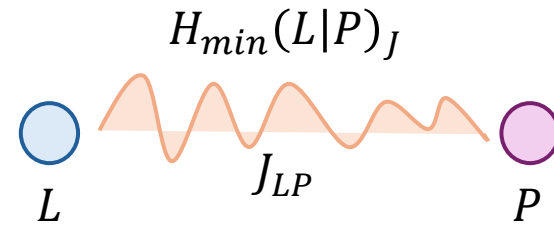
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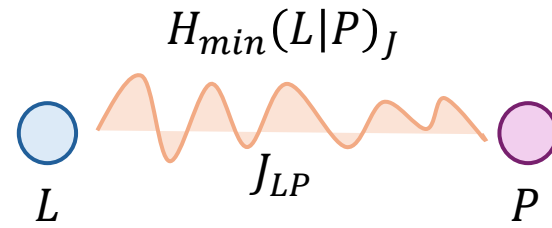


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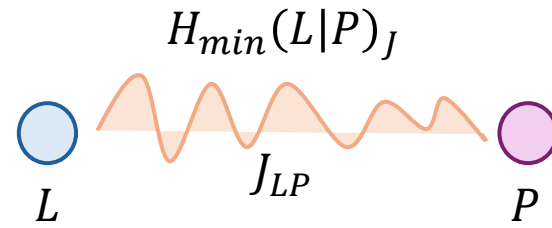
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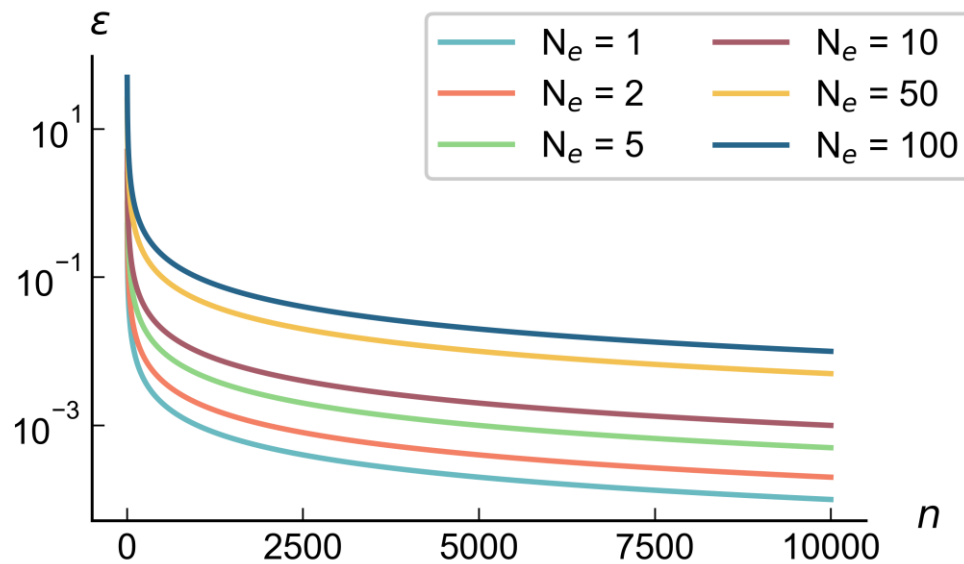
Gives flavour for  $\varepsilon = 0$  no-go result: *erasure noise breaks entanglement!*

$$\Rightarrow \varepsilon \geq \frac{1}{d_L(d_L+1)} \frac{1}{n^2} \quad (n \gg 1)$$

# Example: W-state code

$$|\psi_L\rangle \rightarrow \frac{1}{\sqrt{n}} (|\psi, \perp, \perp, \dots, \perp\rangle + |\perp, \psi, \perp, \dots, \perp\rangle + \dots |\perp, \perp, \perp, \dots, \psi\rangle), \quad \langle\psi|\perp\rangle = 0$$

$$U_L \rightarrow U_L^{\otimes n}$$



The W-state code can  $\varepsilon$ -correct erasure of  $N_e$  subsystems iff

$$\varepsilon \geq \frac{N_e}{n} \left( 1 - \frac{1}{d_L} \right)$$

*e.g.  $n = 100$  qutrits can encode 1 logical qubit and correct for single subsystem erasure up to  $\varepsilon = 0.005$*

# Optimization-free metric

Define “near-optimal infidelity” [1]:

$$\delta(\mathcal{N} \circ \mathcal{E}) := c_L^{-1} (1 - d_L^{-2} \|\text{tr}_L \sqrt{M}\|_2^2)$$

QEC matrix  $M$ :

$$M_{[i,n][j,m]} := \langle i_L | N_n^\dagger N_m | j_L \rangle$$

- $\{N_m\}$  Kraus ops of  $\mathcal{N}$
- $\{|j_L\rangle\}$  code words of  $\mathcal{E}$

(= data needed for Knill-Laflamme)

**Corollary.** Any encoder  $\mathcal{E}_{L \rightarrow P}$  with a universal transversal set of gates can  $\varepsilon$ -correct for erasure of  $m$  subsystems...

$$\text{if: } \varepsilon \geq \delta(\text{tr}_{P_1 \dots P_m} \circ \mathcal{E})$$

$$\text{only if: } \varepsilon \geq \frac{1}{2} \delta(\text{tr}_{P_1 \dots P_m} \circ \mathcal{E})$$

$$\text{cost: } O\left(\left(d_L \prod_{i=1}^m d_{P_i}\right)^3\right) \text{ vs. } O\left((d_L d_P)^{5.246}\right)$$

# Part 3: Methods

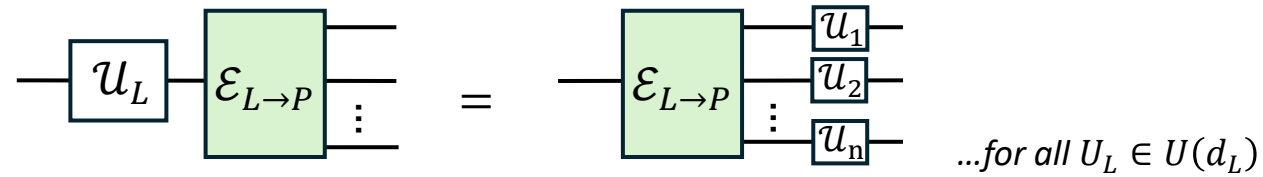
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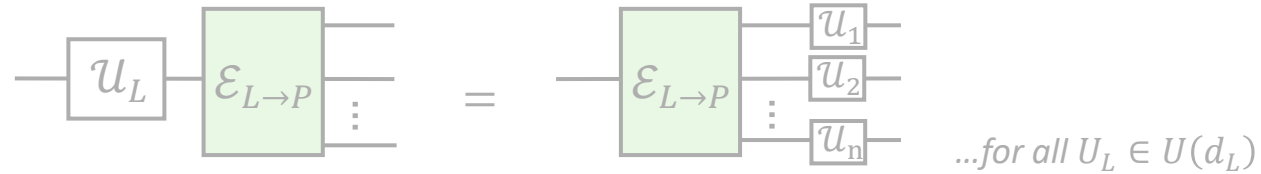
2) Universal transversal gates  $\Rightarrow$  encoder is  $U(d)$ -covariant



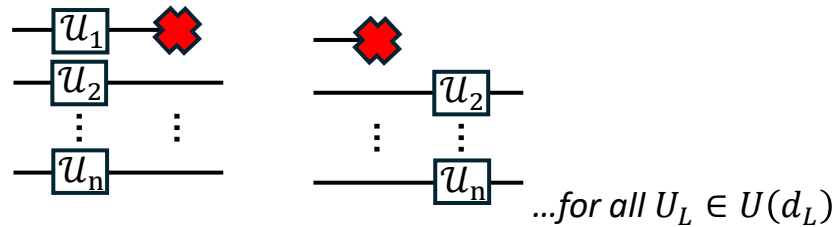
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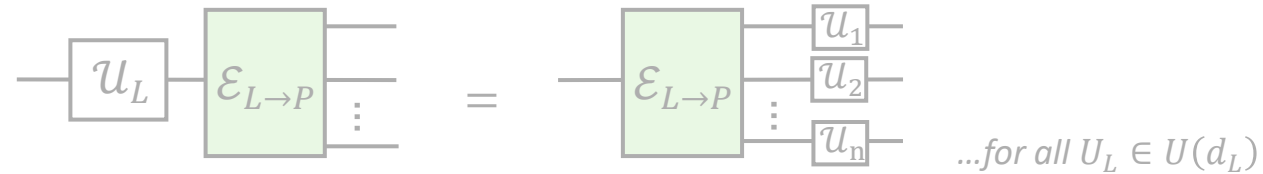
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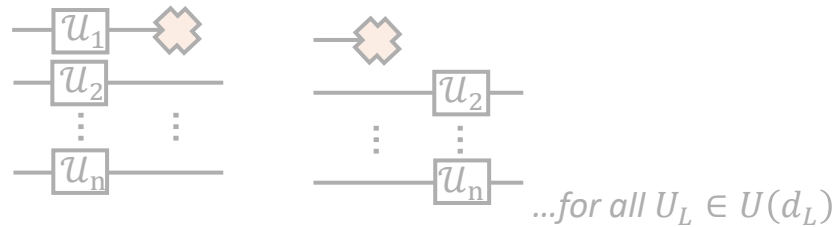
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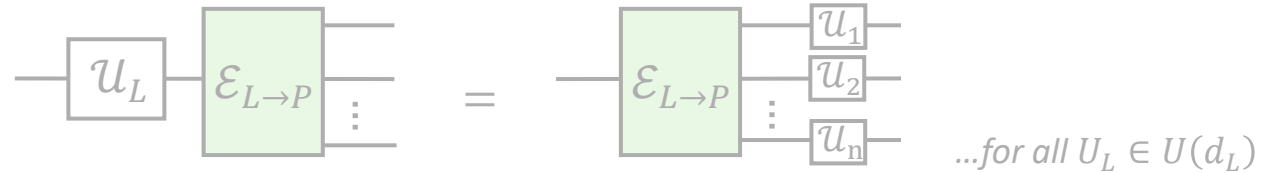
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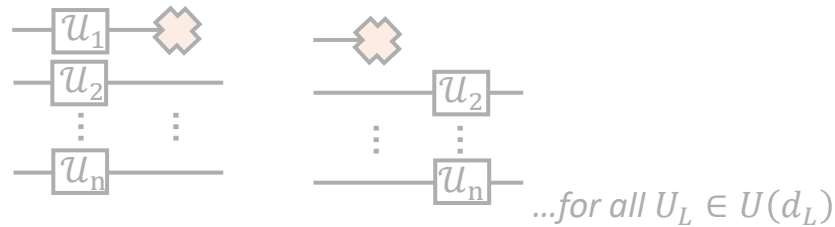
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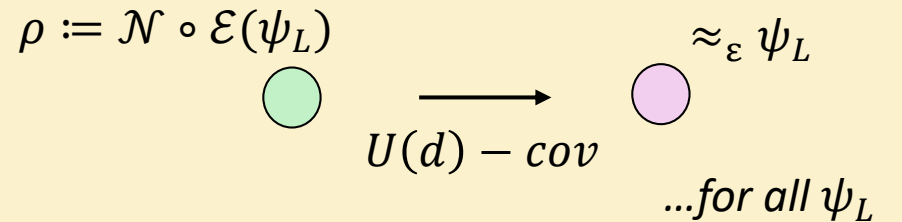
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$\Rightarrow$  QEC as asymmetry distillation task



# Resource theory of asymmetry

★ Free operations = **G-covariant** channels

$$\begin{array}{c} \text{---} \square \varepsilon \text{---} \square \mathcal{U}_g \text{---} \\ \text{---} \square \mathcal{V}_g \text{---} \square \varepsilon \text{---} \end{array} = \quad \forall g \in G$$

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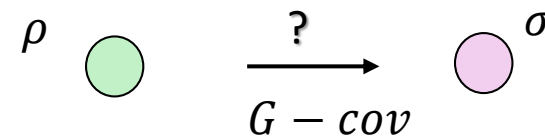
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Q: What can we say about future dynamics of a system in initial configuration  $\rho$  by appealing to symmetries?



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$$\phi \text{ (green circle)} \xrightarrow{\mathcal{V}_G} \psi \text{ (pink circle)} \iff \text{tr}[\phi f(L)] = \text{tr}[\psi f(L)], \forall \text{ continuous } f, \text{ generators } L$$

[Noether 1918]

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**[this work]**  
 \*symmetry + additional linearity constraints\*

# Take home message – future work



(Single-shot) entropic constraint on unitarily-covariant codes

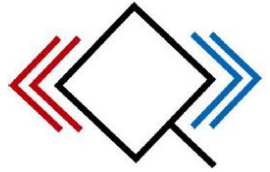


- *Other symmetries? E.g. Clifford group [no-go Chakraborty-Gottesman 2026] / permutation invariant codes...*
- *Noise models beyond erasure [Gupta et al. 2024]?*
- *Explicit construction of efficient (?) optimal decoders for covariant codes*
- *Asymptotic analysis?*
- *Minimal supersets of covariant encoders for circumventing E-K a la [Liu & Zhou 2023]*

★ Thanks!! ★



*Quantum thermodynamics and computing Granada (QTCG)*



Jara Juana  
Bermejo Vega



Daniel  
Manzano



Rhea  
Alexander



Álvaro  
Tejero



Dolores  
Esteve-Díaz



Antonio Jesús  
Rivera Pérez



Carlos  
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Skotiniotis



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